Asymmetric quantum hypothesis testing with Gaussian states

Gaetana Spedalieri and Samuel L. Braunstein

Department of Computer Science, University of York, York YO10 5GH, United Kingdom

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We consider the asymmetric formulation of quantum hypothesis testing, where two quantum hypotheses have different associated costs. In this problem, the aim is to minimize the probability of false negatives and the optimal performance is provided by the quantum Hoeffding bound. After a brief review of these notions, we show how this bound can be simplified for pure states. We then provide a general recipe for its computation in the case of multimode Gaussian states, also showing its connection with other easier-to-compute lower bounds. In particular, we provide analytical formulas and numerical results for important classes of one- and two-mode Gaussian states.

I. INTRODUCTION

Quantum hypothesis testing (QHT) is a fundamental topic in quantum information theory [1,2], playing a nontrivial role in protocols of quantum communication and quantum cryptography [3,4]. The typical formulation of QHT is given in terms of quantum state discrimination [5–8], where the optimal performance is provided by the quantum state discrimination [5–8].

In this paper, we consider asymmetric quantum hypothesis testing, where two quantum hypotheses have different associated costs. In this problem, the aim is to minimize the probability of false negatives and the optimal performance is provided by the quantum Hoeffding bound. After a brief review of these notions, we show how this bound can be simplified for pure states. We then provide a general recipe for its computation in the case of multimode Gaussian states, also showing its connection with other easier-to-compute lower bounds. In particular, we provide analytical formulas and numerical results for important classes of one- and two-mode Gaussian states.

II. BRIEF REVIEW OF ASYMMETRIC TESTING

A. Basic formulation

In binary QHT we consider a quantum system which is prepared in some unknown quantum state $\rho$, which can be $\rho_0$ or $\rho_1$. For instance, we can imagine one party, say Alice, who prepares such a system. This system is then passed to Bob, who does not know which choice Alice has made. Thus, Bob must decide between the following two hypotheses:

Null hypothesis $H_0 : \rho = \rho_0$.

Alternative hypothesis $H_1 : \rho = \rho_1$.

In order to discriminate between these two hypotheses, we consider the quantum state discrimination [5–8], where the optimal performance is provided by the quantum Hoeffding bound. After a brief review of these notions, we show how this bound can be simplified for pure states. We then provide a general recipe for its computation in the case of multimode Gaussian states, also showing its connection with other easier-to-compute lower bounds. In particular, we provide analytical formulas and numerical results for important classes of one- and two-mode Gaussian states.

1. INTRODUCTION

Quantum hypothesis testing (QHT) is a fundamental topic in quantum information theory [1,2], playing a nontrivial role in protocols of quantum communication and quantum cryptography [3,4]. The typical formulation of QHT is given in terms of quantum state discrimination [5–8], where the optimal performance is provided by the quantum state discrimination [5–8].

For such a basic problem, we know closed analytical formulas identifying both the minimum error probability, given by the Helstrom bound [6], and the optimal quantum detection, expressed in terms of the Helstrom matrix [6]. Furthermore, we can also use an easier-to-compute bound which becomes tight in asymptotic conditions. This is the recently introduced quantum Chernoff bound [9], for which we know simple formulas in the case of multimode Gaussian states [10] (i.e., those states with Gaussian Wigner function [5]).

In this paper, we consider asymmetric QHT, where two quantum hypotheses have different associated costs [6–8]. In this approach, we aim to minimize the probability that the alternative hypothesis is confused for the null hypothesis, an error which is known as “false negative.” This minimization has to be done by suitably constraining the probability of another possible error, known as a “false positive,” where the null hypothesis is confused for the alternative hypothesis. This is clearly the best approach, for instance, in medical-type testing, where the null hypothesis typically represents absence of a disease, while the alternative corresponds to the presence of a disease.

Asymmetric QHT is typically formulated as a multicopy discrimination problem, where a large number of copies of the two possible states are prepared and subjected to a collective quantum measurement. From this point of view, the aim is to maximize the error exponent describing the exponential decay of the false negatives, while placing a reasonable constraint on the false positives. For this calculation, we can rely on two mathematical tools. The first is the quantum relative entropy [5] between the two states, while the other is the recently introduced quantum Hoeffding bound (QHB) [11], which performs the optimization of the error exponent while providing a better control on the false positives.

In this work, we start by giving some basic notions on asymmetric QHT and briefly reviewing the QHB, also showing how its computation simply reduces to the quantum fidelity [12] in the presence of pure states. Then, we provide a general recipe for computing this bound in the case of multimode Gaussian states, for which it can be expressed in terms of their first- and second-order statistical moments. In the general multimode case, we derive a relation between the QHB and other easier-to-compute bounds, which are based on well-known mathematical inequalities. Finally, we derive analytical formulas and numerical results for the most important classes of one-mode and two-mode Gaussian states.

By developing the theory of asymmetric QHT for Gaussian states, our work could be useful in tasks and protocols involving Gaussian quantum information [5], including technological applications of quantum channel discrimination (e.g., quantum illumination [13,14] or quantum reading [15–18]).
always reduce his measurement to be a dichotomic POVM \( \{ \Pi_k \} \) with \( k = 0, 1 \) [6]. The outcome \( k = 0 \), with POVM operator \( \Pi_0 \), is associated with the null hypothesis \( H_0 \), while the other outcome \( k = 1 \), with POVM operator \( \Pi_1 = I - \Pi_0 \), is associated with the alternative hypothesis \( H_1 \).

Since the two quantum states \( \rho_0 \) and \( \rho_1 \) are generally nonorthogonal, there is a nonzero error probability to confuse the two hypotheses. We can identify two different types of error: Type-I and type-II errors, with associated conditional error probabilities. By definition, the type-I error, also known as a “false-positive,” is where Bob accepts the alternative hypothesis \( H_1 \) when the null hypothesis \( H_0 \) holds. We have a corresponding error probability expressed by

\[
\alpha := p(H_1|H_0) = \text{Tr}(\Pi_1\rho_0).
\]

(3)

Then, the type-II error or “false negative” is where Bob accepts the null hypothesis \( H_0 \) when the true hypothesis is the alternative \( H_1 \). This error occurs with conditional probability,

\[
\beta := p(H_0|H_1) = \text{Tr}(\Pi_0\rho_1).
\]

(4)

Note that we can introduce other probabilities, but they are fully determined by \( \alpha \) and \( \beta \). For instance, we may also consider the “specificity” or “true negativity” of the test which is the success probability of identifying the null hypothesis, i.e., \( p(H_0|H_0) \), which is simply given by \( 1 - \alpha \). Similarly, we may also consider the “sensitivity” or “true positivity” of the test which is the success probability of identifying the alternative hypothesis, i.e., \( p(H_1|H_1) = 1 - \beta \).

The costs associated with the two types of error can be very different especially in the medical and histological settings. For instance, in a medical test, \( H_0 \) is typically associated with no illness, while \( H_1 \) with the presence of the disease. It is therefore clear that we would like to have tests where the false-negative probability (or rate) \( \beta \) is the lowest possible, so that ill patients are not diagnosed as healthy. For this reason, in a medical setting, hypothesis testing is almost always asymmetric, meaning that we aim to minimize one of the two conditional error probabilities.

### B. Multicopy formulation

In general we can formulate the problem of QHT as an \( M \)-copy discrimination problem [7,8]. This means that Alice has \( M \) quantum systems which are prepared in two possible multicopy states,

\[
H_0 : \rho = \rho_0^\otimes M = \rho_0 \otimes \cdots \otimes \rho_0,
\]

\[
H_1 : \rho = \rho_1^\otimes M = \rho_1 \otimes \cdots \otimes \rho_1.
\]

(5)

These systems are passed to Bob who performs a collective measurement on them. As before, this general POVM can be chosen to be dichotomic \( \{ \Pi_0, \Pi_1 \} \) with \( \Pi_1 = I - \Pi_0 \).

The error probabilities now depend on the number of copies \( M \). In particular, the probability of false positives is given by

\[
\alpha_M := p(H_1|H_0) = \text{Tr}(\Pi_1\rho_0^\otimes M),
\]

(6)

and the probability of false negatives is

\[
\beta_M := p(H_0|H_1) = \text{Tr}(\Pi_0\rho_1^\otimes M).
\]

(7)

In the limit of a large number of copies (\( M \gg 1 \)), these probabilities go to zero exponentially, i.e., we have

\[
\alpha_M \simeq \frac{1}{2} e^{-\alpha_M M}, \quad \beta_M \simeq \frac{1}{2} e^{-\beta_M M},
\]

where the coefficients,

\[
\alpha_R = -\lim_{M \to +\infty} \frac{1}{M} \ln \alpha_M,
\]

(9)

\[
\beta_R = -\lim_{M \to +\infty} \frac{1}{M} \ln \beta_M,
\]

(10)

are called the “error exponents” or “rate limits” [11].

Bob’s aim is to maximize the error exponent \( \alpha_R \), so that the error probability of false negatives \( \beta_M \) has the fastest exponential decay to zero. This must be done while controlling the rate of false positives. Here a well-known result is the “quantum Stein lemma” [11] which connects \( \beta_R \) with the quantum relative entropy between the single-copy states \( \rho_0 \) and \( \rho_1 \). For a large number of copies \( M \gg 1 \), there is a dichotomic POVM such that the error probability of the false positives is bounded,

\[
\alpha_M \leq \varepsilon \quad \text{for any} \quad 0 < \varepsilon < 1,
\]

(11)

and the error probability of false negatives goes to zero with the error exponent,

\[
\beta_R = S(\rho_0||\rho_1) = \text{Tr}\rho_0 (\ln \rho_0 - \ln \rho_1).
\]

(12)

More powerfully, we may use the notion of the QHB [11]. For \( M \gg 1 \), there is a dichotomic POVM such that the error exponent of false positives is lower bounded by a positive parameter,

\[
\alpha_R \geq r \quad \text{for any} \quad r > 0,
\]

(13)

and the error exponent of false negatives satisfies

\[
\beta_R = H(r),
\]

(14)

where \( H(r) \geq 0 \) is the QHB defined by

\[
H(r) := \sup_{0 \leq s < 1} P(r,s), \quad P(r,s) := \frac{-r s - \ln C_s}{1 - s},
\]

(15)

where

\[
C_s := \text{Tr}(\rho_0^s \rho_1^{1-s})
\]

(16)

is the “s overlap” between the single-copy states \( \rho_0 \) and \( \rho_1 \). Note that the quantum Hoeffding bound enforces a stronger constraint on false positives, since these are bounded at the level of the error exponent and not at the level of the error probability as happens for the quantum relative entropy bound.

### III. ASYMMETRIC TESTING WITH PURE STATES

Asymmetric testing becomes very simple when one of the states (or both) is pure. In this case, we can in fact relate the QHB to the quantum fidelity between the two states.

Let us start by considering the case where only one of the states is pure, e.g., \( \rho_0 = |\psi_0\rangle \langle \psi_0| \). We can write [19]

\[
\inf_s C_s = F(|\psi_0\rangle, \rho_1),
\]

(17)

where \( F \) is the fidelity between \( |\psi_0\rangle \) and \( \rho_1 \). Equation (17) implies \( C_s \geq F \). By using the latter inequality in Eq. (15), we
derive the fidelity bound,

$$H(r) \leq H_F(r) := \sup_{0 \leq s < 1} \frac{-rs - \ln F}{1 - s}. \tag{18}$$

This bound can be further simplified by explicitly performing the maximization with regard to the parameter $s$. After a simple calculation we find

$$H_F(r) = \begin{cases} \ln \frac{1}{r}, & \text{for } r \geq \frac{1}{r}, \\ +\infty, & \text{for } r < \frac{1}{r}, \end{cases} \tag{19}$$

which depends on the comparison between the parameter $r$ and the fidelity $F$ of the two states.

More specifically, in the discrimination of two pure states, we find that the previous fidelity bound becomes tight,

$$H(r) = H_F(r). \tag{20}$$

In fact, for pure states $\rho_0 = |\psi_0\rangle\langle\psi_0|$ and $\rho_1 = |\psi_1\rangle\langle\psi_1|$, and for any $0 < s < 1$, we can write

$$C_s = \text{Tr}(|\psi_0\rangle\langle\psi_0|^s|\psi_1\rangle\langle\psi_1|^{1-s}) = \text{Tr}(|\psi_0\rangle\langle\psi_0|^{1-s}|\psi_1\rangle\langle\psi_1|^s)$$

$$= |\langle\psi_0|\psi_1\rangle|^2 = F(|\psi_0\rangle,|\psi_1\rangle). \tag{21}$$

Therefore we can replace $\ln C_s = \ln F$ in the QHB of Eq. (15), which implies Eq. (20) [20].

IV. ASYMMETRIC TESTING WITH GAUSSIAN STATES

A. Basics of bosonic systems and Gaussian states

A bosonic system of $n$ modes is a quantum system described by a tensor product Hilbert space $\mathcal{H}^{\otimes n}$ and a vector of quadrature operators [21,22]:

$$\hat{x}_i^T := (\hat{q}_1, \hat{p}_1, \ldots, \hat{q}_n, \hat{p}_n). \tag{22}$$

These operators satisfy the vectorial commutation relations [23],

$$[\hat{x}_i, \hat{x}_j^T] := [\hat{x}_i^T - (\hat{x}_i^T)^T] = 2i \Omega,$$

where $\Omega$ is the symplectic form, defined as

$$\Omega := \bigoplus_{k=1}^n \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}. \tag{24}$$

Correspondingly, a real matrix $S$ is called “symplectic” when it preserves $\Omega$ by congruence, i.e., $S \Omega S^T = \Omega$.

By definition, we say that a bosonic state $\rho$ is “Gaussian” when its phase-space Wigner representation is Gaussian [5]. In such a case, we can completely describe the state by means of its first- and second-order statistical moments. These are the mean value or displacement vector $\bar{x} := \text{Tr}(\hat{x}_i \rho)$, and the covariance matrix (CM) $V$ with the generic element,

$$V_{ij} = \frac{1}{2} \text{Tr}(\{\hat{x}_i, \hat{x}_j\}_\rho) - \bar{x}_i \bar{x}_j, \tag{25}$$

where $\{,\}$ denotes the anticommutator. The CM is a $2n \times 2n$ real symmetric matrix, which must satisfy the uncertainty principle [5],

$$V + i \Omega \succeq 0. \tag{26}$$

An important tool in the manipulation of Gaussian states is Williamson’s theorem [5]: For any CM $V$, there is a symplectic matrix $S$ such that

$$V = SWS^T. \tag{27}$$

where

$$W = \bigoplus_{k=1}^n v_k \mathbf{I}, \quad \mathbf{I} := \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}. \tag{28}$$

The matrix $W$ is the “Williamson form” of $V$, and the set $\{v_1, \ldots, v_n\}$ is the “symplectic spectrum” of $V$. According to the uncertainty principle, each symplectic eigenvalue must satisfy the condition $v_k \geq 1$, with $v_k = 1$ for all $k$ if and only if the Gaussian state is pure.

B. Computation of the quantum Hoeffding bound

Our goal is to find a general recipe for the calculation of the QHB for Gaussian states. We start from the general formula in Eq. (15) involving the logarithm of the $s$ overlap $C_s$ defined in Eq. (16). Given two $n$-mode Gaussian states, $\rho_0$ and $\rho_1$, we can write an explicit Gaussian formula for the $s$ overlap in terms of their statistical moments ($\bar{x}_0, V_0$) and ($\bar{x}_1, V_1$). This is given by [10,19]

$$C_s = \frac{\Pi_s}{\sqrt{\det \Sigma_s}} \exp \left[ -\frac{d^T \Sigma_s^{-1} d}{2} \right], \tag{29}$$

where $d := \bar{x}_0 - \bar{x}_1$ is the difference between the mean values, while $\Pi_s$ and $\Sigma_s$ depend on the CMs $V_0$ and $V_1$. In particular, introducing the two real functions,

$$G_s(x) := \frac{2^s}{(x + 1)^s - (x - 1)^s}, \tag{30}$$

$$\Lambda_s(x) := \frac{(x + 1)^s + (x - 1)^s}{(x + 1)^s - (x - 1)^s}, \tag{31}$$

we can write the formulas,

$$\Pi_s := 2^n \Pi_{s=1}^n G_s(v_k^0)G_{1-s}(v_k^1), \tag{32}$$

and

$$\Sigma_s := S_0 \left[ \bigoplus_{k=1}^n \Lambda_s(v_k^0)I \right] S_0^T$$

$$+ S_1 \left[ \bigoplus_{k=1}^n \Lambda_{1-s}(v_k^1)I \right] S_1^T, \tag{33}$$

where $\{v_k^0\}$ and $\{v_k^1\}$ are the symplectic spectra of the two states, with $S_0$ and $S_1$ being the symplectic matrices which diagonalize the two CMs according to Williamson’s theorem, i.e.,

$$V_0 = S_0 \left( \bigoplus_{k=1}^n v_k^0 I \right) S_0^T, \quad V_1 = S_1 \left( \bigoplus_{k=1}^n v_k^1 I \right) S_1^T. \tag{34}$$

Substituting Eq. (29) into Eq. (15) corresponds to explicitly computing the logarithmic term $\ln C_s$, yielding

$$\ln C_s = \ln \Pi_s - \frac{d^T \Sigma_s^{-1} d}{2}. \tag{35}$$

In particular for zero-mean Gaussian states we have $d = 0$ and the previous expression simplifies to

$$\ln C_s = \ln \Pi_s - \frac{1}{2} \ln \det \Sigma_s. \tag{36}$$
C. Other computable bounds

Note that computing the $s$ overlap $C_s$ and its logarithmic form $\ln C_s$ could be difficult due to the presence of the symplectic matrices, $S_0$ and $S_1$, in the term $\Sigma_s$ in Eq. (33). A possible solution is to compute an upper bound, known as the “Minkowski bound,” which is based on the Minkowski determinant inequality [24] and depends only on the two symplectic spectra

\[ M_s := 4^n \left[ \prod_{k=1}^{n} \Psi_s(v_k^0, v_k^1) + \prod_{k=1}^{n} \Psi_{1-s}(v_k^0, v_k^1) \right]^{-n}, \]  

and

\[ \Psi_s(x, y) := \frac{[(x + 1)^s + (x - 1)^s]}{[(y + 1)^{1-s} - (y - 1)^{1-s}]}^{1/n}. \]

Another easy-to-compute upper bound is the “Young bound” $Y_s$, which is based on Young’s inequality [25] and satisfies

\[ C_s \leq M_s \leq Y_s, \]  

where [10]

\[ Y_s := 2^n \prod_{k=1}^{n} \Gamma_s(v_k^0) \Gamma_{1-s}(v_k^1), \]  

and

\[ \Gamma_s(x) := \frac{[(x + 1)^s - (x - 1)^s]}{2}^{\frac{1}{s}}. \]

Taking the negative logarithm of Eq. (39), we can write the following inequality for the QHB:

\[ H(r) \geq H_M(r) \geq H_T(r), \]

where

\[ H_M(r) := \sup_{0 \leq s < 1} \frac{-r s - \ln M_s}{1 - s}, \]

\[ H_T(r) := \sup_{0 \leq s < 1} \frac{-r s - \ln Y_s}{1 - s}. \]

In the specific case where one of the two Gaussian states is pure, we can compute their fidelity $F$ and apply the upper bound given in Eqs. (18) and (19), which becomes tight when both states are pure [see Eq. (20)]. In particular, for two multimode Gaussian states $\rho_0 = |\psi_0\rangle \langle \psi_0|$ and $\rho_1$, we can easily write their fidelity $F$ in terms of the statistical moments [19]:

\[ F = \frac{2^n}{\sqrt{\det L}} \exp \left( -\frac{\text{d}^T L^{-1} \text{d}}{2} \right), \]

where $L := V_0 + V_1$. As a result, we can use Eq. (19) with

\[ \ln \frac{1}{F} = \frac{1}{2} \left[ \ln \left( \frac{\det L}{4^n} \right) + \text{d}^T L^{-1} \text{d} \right]. \]

V. DISCRIMINATION OF ONE-MODE GAUSSIAN STATES

In this section, we examine the case of one-mode Gaussian states. This means we fix $n = 1$ in the previous formulas of Sec. IV, with matrices becoming $2 \times 2$, vectors becoming two-dimensional, and symplectic spectra reducing to a single eigenvalue. For instance, the $s$ overlap can be more simply computed using the expressions,

\[ \Pi_s = 2G_s(v^0)G_{1-s}(v^1), \]

\[ \Sigma_s = \Lambda_s(v^0)S_0 + \Lambda_{1-s}(v^1)S_1. \]

In particular, here we shall derive the analytic formulas for the QHB for two important classes: coherent states (in Sec. V A) and thermal states (in Sec. V B).

A. Asymmetric testing of coherent amplitudes

The expression of the QHB is greatly simplified in the case of one-mode coherent states $\rho_0 = |\alpha_0\rangle \langle \alpha_0|$ and $\rho_1 = |\alpha_1\rangle \langle \alpha_1|$. Since both states are pure, the QHB is equal to the fidelity bound in Eq. (19), i.e., $H(r) = H_F(r)$. Therefore, it is sufficient to compute the fidelity between the two coherent states, which is given by

\[ F = |\langle \alpha_0 | \alpha_1 \rangle|^2 = e^{-|\alpha_0 - \alpha_1|^2}, \]

so that $\ln \frac{1}{F} = |\alpha_0 - \alpha_1|^2 := \sigma$, and we can write

\[ H(r) = \frac{\sigma}{r}, \quad \text{for } r \geq \sigma, \]

\[ +\infty, \quad \text{for } r < \sigma. \]

Assuming that we impose a good control on the rate of false positives (so that $r \geq \sigma$), then the error exponent for the false negatives is simply given by $H(r) = \sigma$. More explicitly, this corresponds to an asymptotic error rate,

\[ \beta_M = \frac{1}{2} e^{-M\sigma} = \frac{e^M}{2}. \]

Note that if we have poor control on the rate of false positives, i.e., $r < \sigma$, then the QHB $H(r)$ is infinite. This means that the probability of false negatives $\beta_M$ goes to zero superexponentially, i.e., more quickly than any decreasing exponential function.

B. Asymmetric testing of thermal noise

In this section we derive the QHB for one-mode thermal states $\rho_0 = \rho_{\text{BH}}(v^0)$ and $\rho_1 = \rho_{\text{BH}}(v^1)$, with variances equal to $v^0$ and $v^1$, respectively (in our notation, $v = 2n + 1$, where $n$ is the mean number of thermal photons). These Gaussian states have zero mean ($\bar{x}_0 = \bar{x}_1 = 0$) and CMs in the Williamson form $V_0 = v^0 I$ and $V_1 = v^1 I$ (so that $S_0 = S_1 = I$). Thus, we can write

\[ \Sigma_s = \epsilon_s I, \quad \epsilon_s := \Lambda_s(v^0) + \Lambda_{1-s}(v^1), \]

and derive

\[ C_s = \frac{\Pi_s}{\epsilon_s} = \frac{2}{(v^0 + 1)^{v^1 + 1} - (v^0 - 1)^{v^1 + 1}}. \]

This is the $s$ overlap to be used in the QHB of Eq. (15).

Given two arbitrary $v^0 \geq 1$ and $v^1 \geq 1$, the maximization in Eq. (15) can be done numerically. The results are shown in Fig. 1 for thermal states with variances up to 3 vacuum units (equivalent to 1 mean thermal photon). From the figure we see an asymmetry with respect to the bisector $v^0 = v^1$ which is a consequence of the asymmetric nature of the hypothesis test.
Let us now consider the thermal state to be the null hypothesis ($ν^0 := ν > 1$) while the vacuum state is the alternative hypothesis ($ν^1 = 1$). In this case, we derive

$$ P(r,s) = \frac{s}{1-s} \left[ \ln \left( \frac{1 + ν^1}{2} \right) - r \right], \quad (57) $$

which leads to the following expression for the QHB:

$$ H(r) = \begin{cases} 0 & \text{for } r \geq \ln \left( \frac{1+ν^1}{2} \right), \\ +\infty & \text{for } r < \ln \left( \frac{1+ν^1}{2} \right). \end{cases} \quad (58) $$

This is related to the minimum probability of confusing the vacuum state with a thermal state. Note that this is very different from Eq. (56).

VI. DISCRIMINATION OF TWO-MODE GAUSSIAN STATES

In this section we consider two important classes of two-mode Gaussian states. The first is the class of Einstein-Podolsky-Rosen (EPR) states, also known as two-mode squeezed vacuum states. The second (broader) class is that of two-mode squeezed thermal (ST) states, for which the computation of the QHB is numerical.

A. Asymmetric testing of EPR correlations

The expression of the QHB in the case of EPR states is easy to derive. Since EPR states are pure, the QHB $H(r)$ is given by $H_F(r)$ of Eq. (19). As a result, we need only to compute the fidelity between the two states.

An EPR state has zero mean and CM,

$$ V_{\text{EPR}}(μ) = \begin{pmatrix} μ & \sqrt{μ^2 - 1}Z \\ \sqrt{μ^2 - 1} & μ \end{pmatrix}, \quad (59) $$

with $μ \geq 1$, $I$ is the $2 \times 2$ identity matrix, and

$$ Z := \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}. \quad (60) $$

Given two EPR states with parameters $μ_0$ and $μ_1$, their fidelity is computed via Eq. (45), yielding

$$ F = \frac{4}{\sqrt{\det L}}, \quad (61) $$

where $L = V_{\text{EPR}}(μ_0) + V_{\text{EPR}}(μ_1)$. After simple algebra, we find

$$ F = \frac{2}{\left( 1 + μ_0μ_1 - (μ_0^2 - 1)(μ_1^2 - 1) \right)}, \quad (62) $$

to be used in Eq. (19).

B. Squeezed thermal states

In this section we consider symmetric ST states $ρ(μ,c)$, which are Gaussian states with zero mean and CM,

$$ V_{\text{ST}}(μ,c) = \begin{pmatrix} μ & cZ \\ cZ & μ \end{pmatrix}, \quad (63) $$

where $μ \geq 1$ and $|c| \leq μ$ [26,27] (in particular, without loss of generality, we can assume $c \geq 0$). These are called symmetric
The symplectic matrix $S$ which diagonalizes $V_{ST}(\mu,c)$ in the Williamson form $v(I \oplus I)$ is given by

$$S = \begin{pmatrix} \omega_+ & 1 \\ \omega_- Z & \omega_- I \end{pmatrix},$$

where

$$\omega_\pm := \frac{\sqrt{\mu \pm v}}{2v}.$$  

As a result, the overlap between two symmetric ST states, $\rho_0$ and $\rho_1$, can be computed using the simplified formulas,

$$\Pi_s = 4G_0^2(v^0)G_{1-s}(v^1),$$

$$\Sigma_s = \Lambda_s(v^0)S_0T_S + \Lambda_{1-s}(v^1)S_1T_S,$$

where $v^0(v^1)$ is the degenerate eigenvalue of $\rho_0(\rho_1)$, computed according to Eq. (64), and $S_0(S_1)$ is the corresponding diagonalizing symplectic matrix, computed according to Eqs. (65) and (66).

Let us start with simple cases involving the asymmetric testing of correlations with specific ST states. First we consider the asymmetric discrimination between the uncorrelated thermal state $\rho_0 = \rho(\mu,0)$ as null hypothesis and the correlated (but separable) ST state $\rho_1 = \rho_{EPR}(\mu)$ as alternative hypothesis. A false negative corresponds to concluding that there are no correlations where they are actually present [29]. It is straightforward to derive their degenerate symplectic eigenvalues which are simply $v^0 = \mu$ and $v^1 = \sqrt{2\mu - 1}$. Then, we have $S_0 = I \oplus I$, while $S_1$ can be easily computed from Eqs. (65) and (66). By substituting these into Eqs. (67) and (68), we can compute the overlap $C_s = \Pi_s/\sqrt{\det \Sigma_s}$, and therefore the QHB $H(r)$ via Eq. (15). The results are plotted in Fig. 2, for values of thermal variance $\mu$ up to 3 (i.e., from zero to 1 mean photon) and small values of the parameter $r$, bounding the rate of false positives. As expected, the QHB improves for decreasing $r$ and increasing $\mu$.

Now let us consider the asymmetric discrimination between $\rho_0 = \rho(\mu,0)$ and the EPR state $\rho_{EPR}(\mu)$, i.e., the most correlated and entangled ST state [29]. Thanks to the simple symplectic decomposition of the EPR state ($v^1 = 1$), we can further simplify the previous Eqs. (67) and (68) and write

$$\Pi_1 = 4G_1^2(\mu), \quad \Sigma_1 = \Lambda_s(\mu)(I \oplus I) + V_{EPR}(\mu).$$

FIG. 2. (Color online) Asymmetric discrimination between the thermal state $\rho_0 = \rho(\mu,0)$ and the ST state $\rho_1 = \rho(\mu,\mu - 1)$ with maximal separable correlations. We plot the QHB as a function of the thermal variance $\mu$ and the false-positive parameter $r$. As we can see the QHB improves for lower $r$ and for higher $\mu$.

FIG. 3. (Color online) Asymmetric discrimination between the thermal state $\rho_0 = \rho(\mu,0)$ and the EPR state $\rho_{EPR}(\mu)$. We plot the QHB as a function of the thermal variance $\mu$ and the false-positive parameter $r$. The QHB improves for lower $r$ and for higher $\mu$. In particular, there is a threshold value after which the QHB becomes infinite (white region).
Finally, we consider the most general scenario in the asymmetric testing of correlations with ST states. In fact, we consider two generic ST states, $\rho(\mu, c_0)$ and $\rho(\mu, c_1)$, with the same thermal noise but differing amounts of correlation. For this computation, we use Eqs. (64)–(66) with $c = c_0$ or $c_1$, to be replaced in Eqs. (67) and (68), therefore deriving the $s$ overlap and the QHB. At small thermal variance ($\mu = 3$) and for the numerical value $r = 0.1$, we plot the QHB as a function of the correlation parameters $c_0$ and $c_1$. As we can see from Fig. 4, the QHB is not symmetric with respect to the bisector $c_0 = c_1$ (where it is zero) and increases away from this line.

![Graph showing asymmetric discrimination between two ST states with the same thermal variance ($\mu = 3$) but different correlations $c_0$ and $c_1$. Setting $r = 0.1$, we plot the QHB as a function of $c_0$ and $c_1$. We can see that the QHB increases orthogonally to the bisector $c_0 = c_1$. The QHB is finite, and the other where it is infinite (white region in the figure).

In fact, by expanding the term $P(r,s)$ in Eq. (15) for $s \to 1^-$, then we find

$$P(r,s) \simeq \frac{N}{s - 1} + O(s - 1),$$

where

$$N := r - \ln \left( \frac{1 + 3\mu^2}{4} \right).$$

For values of $r$ and $\mu$ such that $N > 0$, we find that the term $P(r,s)$ diverges at the border, making the QHB infinite. For a given $r$, this happens when

$$\mu > \bar{\mu}(r) := \sqrt{\frac{4e^r - 1}{3}}.$$
We can write \( H(r) = \max\{P(r,0), \sup_{0<s<1} P(r,s)\} \), where \( P(r,0) = -\ln C_0 = 0 \) can be neglected and

\[
\sup_{0<s<1} P(r,s) = \sup_{0<s<1} \frac{-rs - \ln F}{1-s} = \begin{cases} 
\ln \frac{1}{F} & \text{for } r \geq \ln \frac{1}{F}, \\
+\infty & \text{for } r < \ln \frac{1}{F}.
\end{cases}
\]