

# Probabilities and Mixed-Criticalities: the Probabilistic C-Space

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# Why considering a probabilistic analysis ?

- 1 Probabilistic C-Space schedulability analysis
- 2 Mixed-criticality with probabilities: define the probability of being in a given criticality level
- 3 Probabilistic Sensitivity Analysis

**In this presentation we explore Item (2) with the help of the probabilistic C-Space**

# Motivation

Predicting Task Execution Time is difficult:

- execution variability (inputs, environment, tasks)
- additional (and unpredictable) delays in task executions from the system
- interferences from system elements (included other tasks)

pWCET to model task execution variability is an alternative to Worst-Case Execution Times (WCETs)

# Contributions

## Provide probabilistic guarantees for MC

- **formalize** scheduling problem with criticalities through probabilities
- **pC-Space**: pWCETs applied to construct probabilistic version of C-space
- provide **initial evaluation** flexibility of probabilistic models and probabilistic scheduling for the mixed-criticality problem

# Outline

1 Introduction

**2 Modeling**

3 Schedulability and C-space with probabilities

# Computational Model

## Task model and Scheduling

$$\tau_i = (C_i, T_i, D_i)$$

Set of  $n$  periodic tasks  $\Gamma = \{\tau_1, \dots, \tau_n\}$

$$H = \text{lcm}(T_1, \dots, T_n)$$

$\Gamma$  is scheduled with EDF

# Probabilistic Worst-Case Execution Time

The  $\text{pdf}_{C_i}$  is the probability distribution function (pdf) representation of the random variable  $C_i$ .

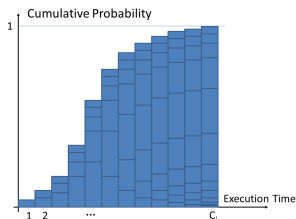
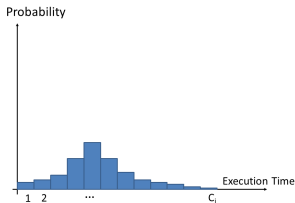
$$\text{pdf}_{C_i} = \left( \begin{array}{ccc} C_{i,1} & \dots & C_{i,k_i} \\ P(C_i = c_{i,1}) & \dots & P(C_i = c_{i,k_i}) \end{array} \right)$$

- $\text{pdf}_{C_i}(c_{i,r}) = P(C_i = c_{i,r})$ ,  $\sum_{r=1}^{k_i} \text{pdf}_{C_i}(c_{i,r}) = 1$ ;  $P(C_i \leq C_i) = 1$

# Probabilistic Worst-Case Execution Time

Cumulative distribution function (cdf):

$$\text{cdf}_{C_i}(c) = P(C_i \leq c) = \sum_{x=1}^c \text{pdf}_{C_i}(x)$$



(a) Worst-case execution time histogram representation (b) Worst-case execution time cdfs



# Probabilistic Worst-Case Execution Time

## WCET Thresholds

$C_i \rightarrow$  WCET thresholds:  $\langle c_{i,k}, p_{i,k} \rangle, 1 \leq k \leq k_i$

The worst-case value  $c_{i,k}$  associated to  $p_{i,k}$

$p_{i,k} \stackrel{\text{def}}{=} \text{cdf}_{C_i}(c_{i,k})$ : **accuracy/confidence** of WCET  $c_{i,k}$

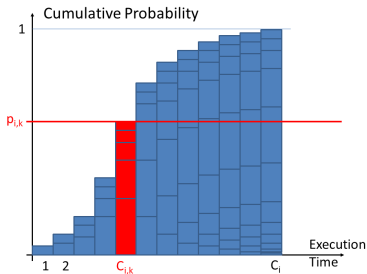


Figure: WCET threshold

# Probabilistic Mixed-Criticality

Classical Mixed-critical WCET  $C_i = (LO - WCET, HI - WCET)$

LO-WCET by a less pessimistic timing analysis tools, larger

HI-WCET by more conservative timing analysis tools

## MC analysis

- Most pessimistic assumption: If one HI-criticality task executes beyond its LO-WCET the system considers that all HI-criticality will do the same
- Our approach: define for each task the level of confidence on its Execution Time w.r.t a criticality level

The probability of meeting deadlines is function of the criticality level that depends on the level of confidence on the WCETs

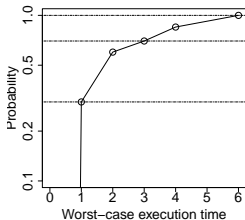
# Probability Threshold

- Design parameter  $\beta$  (**probability threshold**) defining the level of confidence for a WCET imposed to a task at a criticality level *crit*
- From  $\beta$  we get  $c_i(\beta)$  the corresponding WCET in the cdf of  $C_i$
- $\beta \times 100\%$  of the worst-case execution time experienced by  $\tau_i$  are below  $c_i(\beta)$

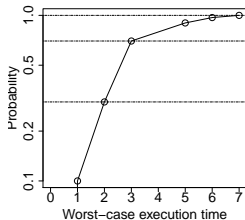
Now for any level of confidence  $\beta$  what is the probability to meet deadlines at criticality level *crit* ?

# Probability Threshold

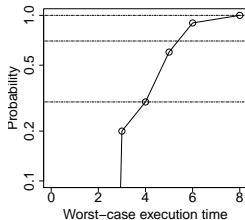
We can define a level of confidence for each criticality level and for each task



(a)  $\tau_1$



(b)  $\tau_2$



(c)  $\tau_3$

# Outline

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# Probabilistic Demand Bound Function

$\text{dbf}_{i,j}(t)$ : the demand bound function in the interval  $[0, t]$ :

$$\text{dbf}_{i,j}(t) \stackrel{\text{def}}{=} \lfloor \frac{t - D_i}{T_i} + 1 \rfloor \times c_{i,j}$$

- $\text{dbf}_{i,k}$  from  $c_{i,k} \leq c_{i,j}$ ;  $\langle \text{dbf}_{i,j}(t), p_{i,j} \rangle$  with  $p_{i,j}$  the probability that  $\text{dbf}_{i,j}(t)$  upper-bounds  $\tau_i$  resource demand in  $[0, t]$  (not function of  $t$ , obtained from the cumulative distribution of  $C_i$ )
- All  $\langle \text{dbf}_i(t, c_{i,j}), p_{i,j} \rangle$  form a distribution of demand bound functions
- $DBF_i(t) = (\text{dbf}_i(t, \cdot), p_i(\cdot))$ : probabilistic demand bound function (probabilistic demand curve) of  $\tau_i$  in  $[0, t]$

# Probabilistic Demand Bound Function

From each  $DBF_i(t)$ , we can then compute the application level probabilistic distribution  $DBF = (dbf(t, \cdot), \rho(\cdot))$ :

$$DBF(t) = \otimes_i DBF_i(t),$$

The EDF schedulability condition states:

$$\forall t \in D, \quad dbf(t) \leq t,$$

with  $D$  the set of deadlines within the hyperperiod (this number can be significantly with a linear programming approach (George & Hermant 2009 for constrained deadlines)

# Probabilistic Demand Bound Function

Now what is the probability to satisfy at time  $t$ ,  $dbf(t) \leq t$  ?

- This can be computed by taking in  $DBF(t)$  the maximum value less than or equal to  $t$  to find the associated probability  $P(t)$

Hence for all times  $t_i$  to consider in  $[0, H]$ , the probability  $P$  to meet meet deadlines with EDF is:

$$P = P(t_1) \times P(t_2) \times \dots$$



# Probability to meet deadlines at level of confidence $\beta$

- A criticality level is defined with its associated level of confidence  $\beta$
- From  $\beta$ , we constrain execution times of each task to a maximum value
- We therefore update  $DBF_i(t)$  to take into account only possible execution times

$$P_\beta = P_\beta(t_1) \times P_\beta(t_2) \times \dots$$

# Probabilistic Schedulability

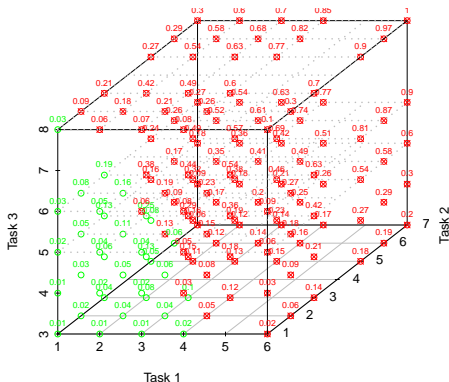
## Theorem ( $\beta$ -EDF Schedulability)

*Any probabilistic task set  $\{\tau_1, \tau_2, \dots, \tau_n\}$  with in a dbf distribution  $DBF(t_i)$  at time  $t_i$  is schedulable with a probability  $P_\beta$  under EDF with a level of confidence  $\beta \in ]0, 1]$ , with  $S$  the set of significant deadlines in  $[0, H]$  and :*

$$P_\beta = P_\beta(t_1) \times P_\beta(t_2) \times \dots$$

# Probabilistic C-Space [1/2]

- We can extend the  $\beta$ -EDF Schedulability condition to characterize the pC-space
- A vector of execution times  $\bar{c} = \{c_1, c_2, \dots\}$  in the pc-Space can be associated to a corresponding probability in all distributions  $DBF(t_i), t_i \in S$



# Probabilistic C-Space [2/2]

The probabilities within the pC-Space can be interpreted in various ways:

- 1 as the confidence of not passing the WCET thresholds of  $\bar{c}$ ;
  - $p(crit) = p_1(crit) \times p_2(crit) \times \dots$
  - $1 - p(crit)$  the possibility of changing that level
- 2 the confidence on system schedulability  $P$ : points inside the region are schedulable, confidence of at least  $P$
- 3 the confidence  $\beta$  on task worst-case behavior;  $\beta$  is the probability of passing  $\bar{c}(\beta)$

# Conclusion and Future Work

Combining probabilities and mixed-criticality problem through:

- probabilistic C-space
- probabilistic sensitivity analysis

Formalization and initial ideas for helping designing and applying MC scheduling policies

# Conclusion and Future Work

In the future:

- define probabilistic sensitivity analysis in terms of change strategies and effect evaluation
- leverage the information from the probabilistic models (pWCET distributions and confidences)
- provide system design feedbacks for an optimal (at least suboptimal) system resource utilization for different criticalities

Thank you!

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