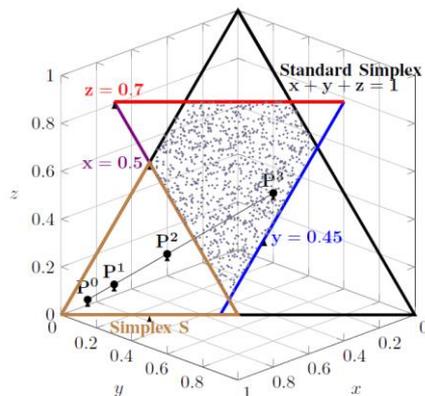


Generating Utilization Vectors for the Systematic Evaluation of Schedulability Tests

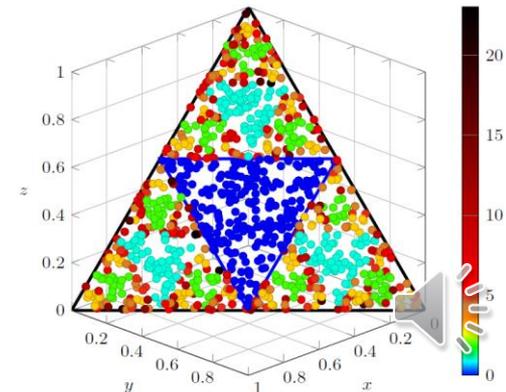


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Real-Time Systems Research Group, University of York, UK



RTSS 2020



Introducing the DRS Algorithm

- **Dirichlet-Rescale (DRS) algorithm**

$$\mathbf{u} = \text{DRS}(n, U, \mathbf{u}^{\max}, \mathbf{u}^{\min})$$

- **Returns:**

A vector of n components (utilization values) $\mathbf{u} = (U_1, U_2, \dots, U_n)$

such that $\sum_{i=1}^n U_i = U$

$$\forall i \ U_i^{\max} \geq U_i \geq U_i^{\min} \geq 0$$

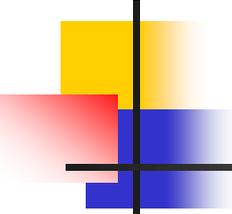
- **Inputs:**

n – size of the vector required

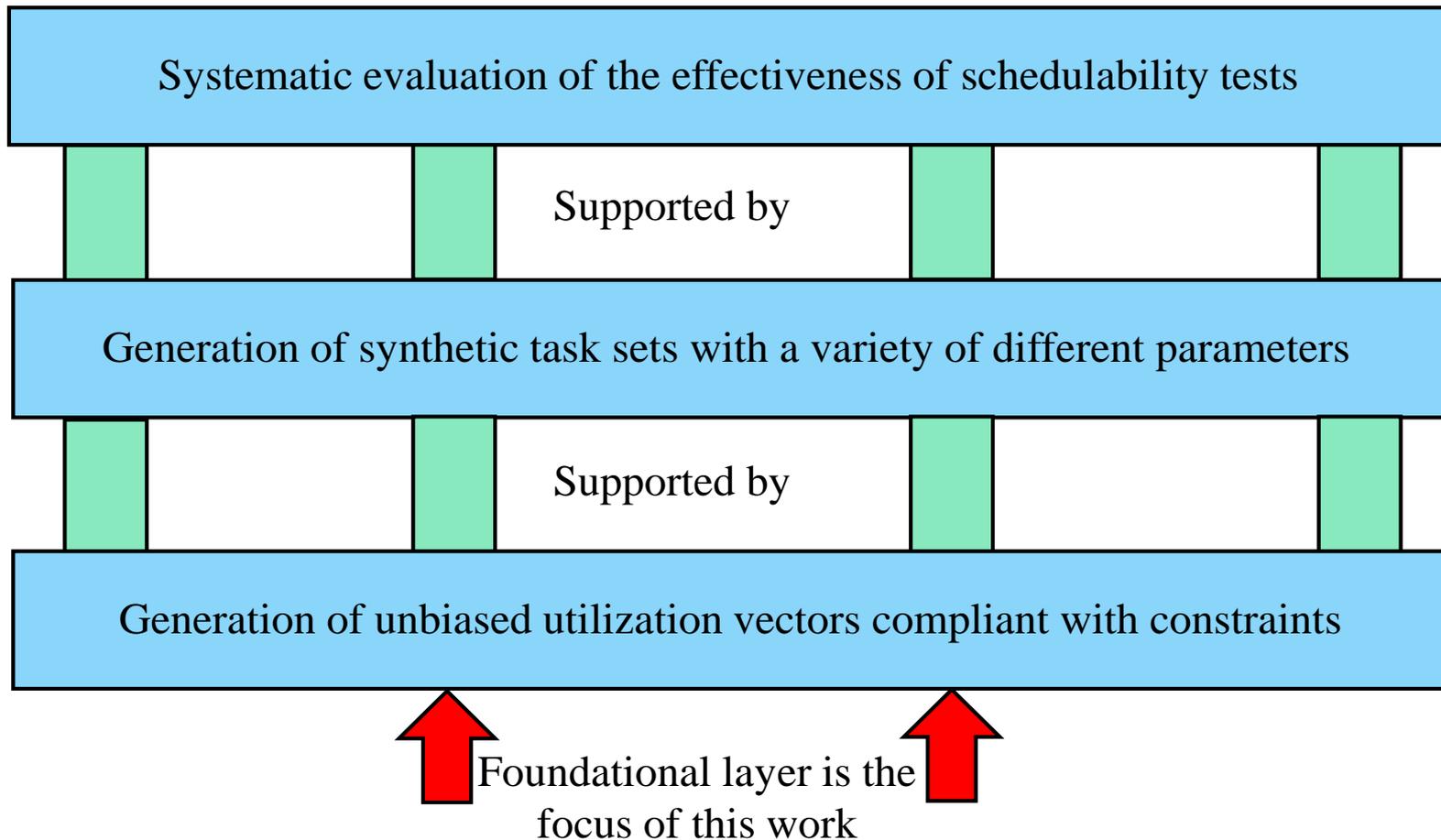
U – total utilization required

$\mathbf{u}^{\max} = (U_1^{\max}, U_2^{\max}, \dots, U_n^{\max})$ vector of maximum constraints

$\mathbf{u}^{\min} = (U_1^{\min}, U_2^{\min}, \dots, U_n^{\min})$ vector of minimum constraints



Motivation



Key criteria for utilization vector generation

■ **Uniformity**

- The vectors of utilization values generated must be unbiased i.e. the vectors must be uniformly distributed within the valid region
 - Bias in the sets of vectors generated can undermine the conclusions drawn from studies into schedulability test effectiveness (Bini and Buttazzo, 2005 [6])

■ **Efficiency**

- Necessary to generate millions of task sets to achieve statistically significant sample sizes in wide-ranging systematic evaluations
 - Typically 1000 task sets per data point for high quality results (Davis, 2016 [11])

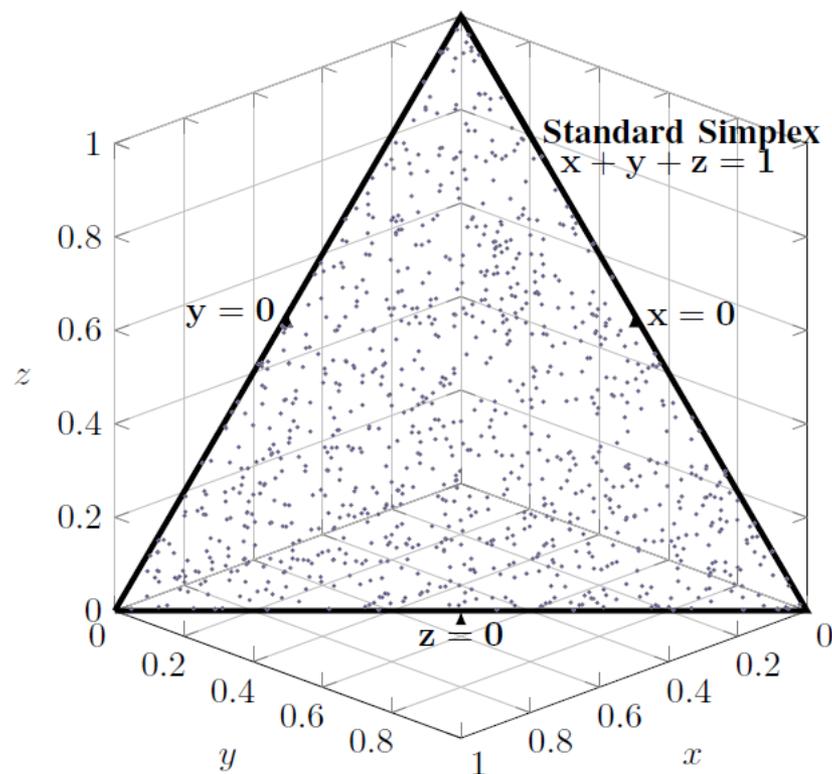
■ **Flexibility**

- Capable of handling constraints on individual task utilization values
 - So the utilization vectors can be tailored to the specific requirements of the problem at hand (examples later), while still producing a uniform distribution of vectors within the valid region given by the constraints

Mathematical background

■ Vectors and Simplices

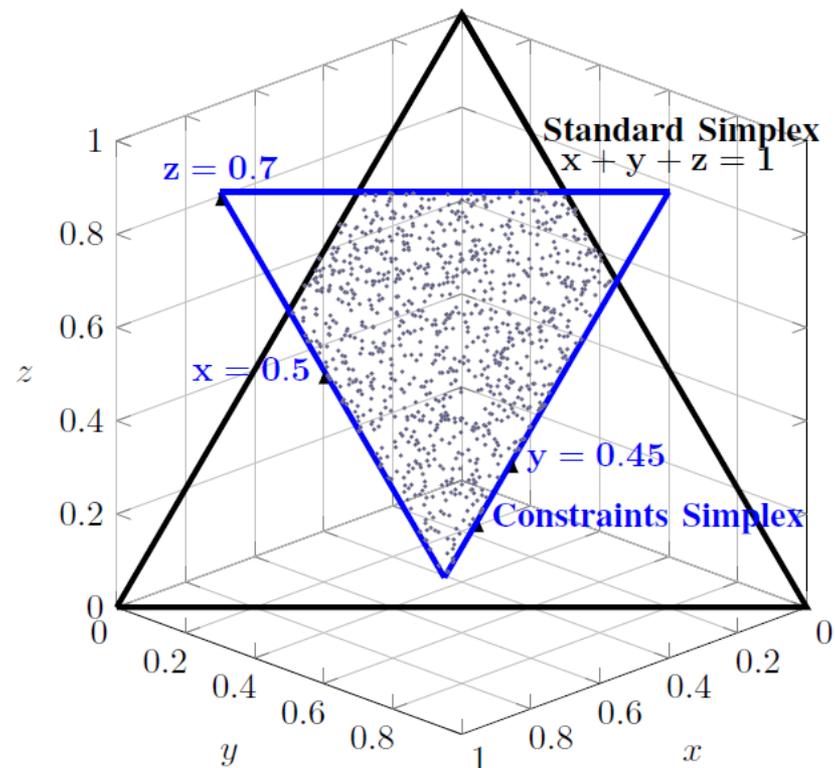
- n-dimensional vectors with components that sum to U
- Each vector represents a point in n-dimensional space ($n=3$ for visualization)
- Canonical form $x+y+z=1$ with $x \geq 0, y \geq 0, z \geq 0$
- Equation $x+y+z=1$ defines a hyperplane (plane in 3-D space)
- Combined with inequalities defines a standard n-1 dimensional simplex embedded in n dimensional space (triangle in 3-D space)
- Vectors required are points uniformly distributed within this simplex



Mathematical background

■ Adding constraints

- Maximum constraints form a *constraints simplex* on the same hyperplane as the standard simplex ($x+y+z=1$ and $x \leq 0.5, y \leq 0.45, z \leq 0.7$)
- Vectors required are points uniformly distributed within the *valid region* i.e. within the intersection of the constraints simplex and the standard simplex
- *Duality* between the two simplices – we could generate points in either simplex and use the other as the constraints
- Minimum constraints can be handled by transforming the problem into a canonical form where all minimum constraints are zero (see the paper)



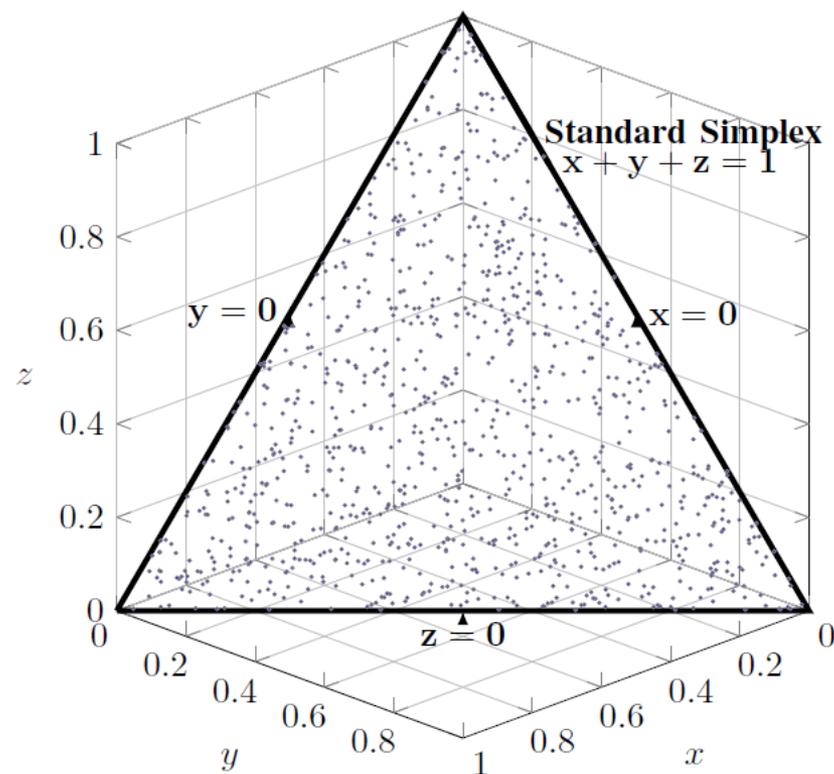
Related work

■ UUnifast algorithm

- First work on this topic published in the Real Time Systems literature
- Bini and Buttazzo, 2005 [6]
- Solves the problem with no maximum or minimum constraints
- Useful for single processor systems

■ Flat Dirichlet distribution

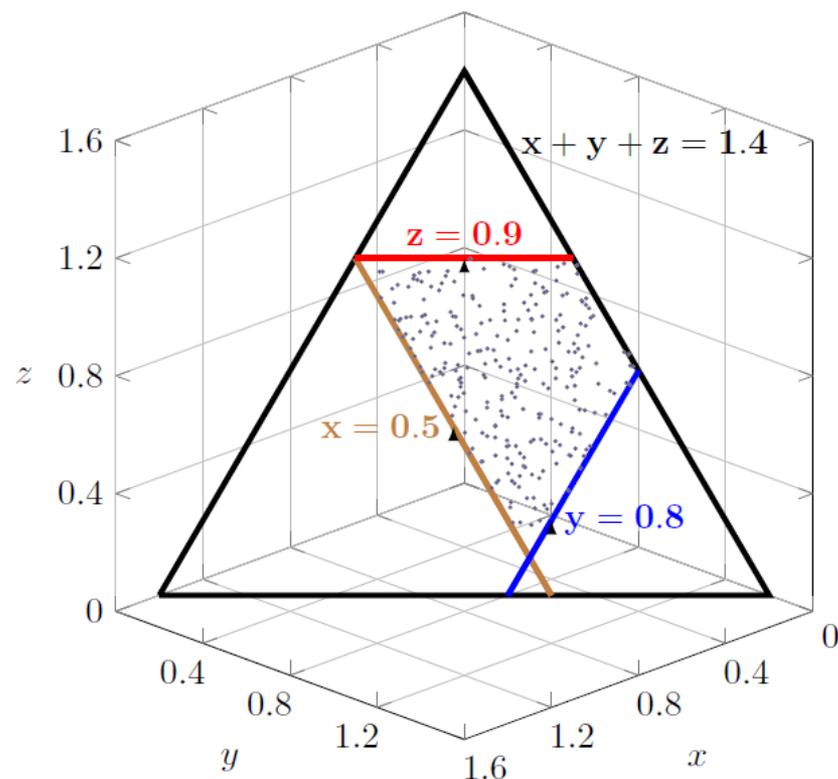
- In the Maths literature, Olkin and Rubin, 1964 [24] published work on the Dirichlet distribution
- Can also be used to solve the problem with no constraints for single processor systems



Related work

■ UUnifast-Discard algorithm

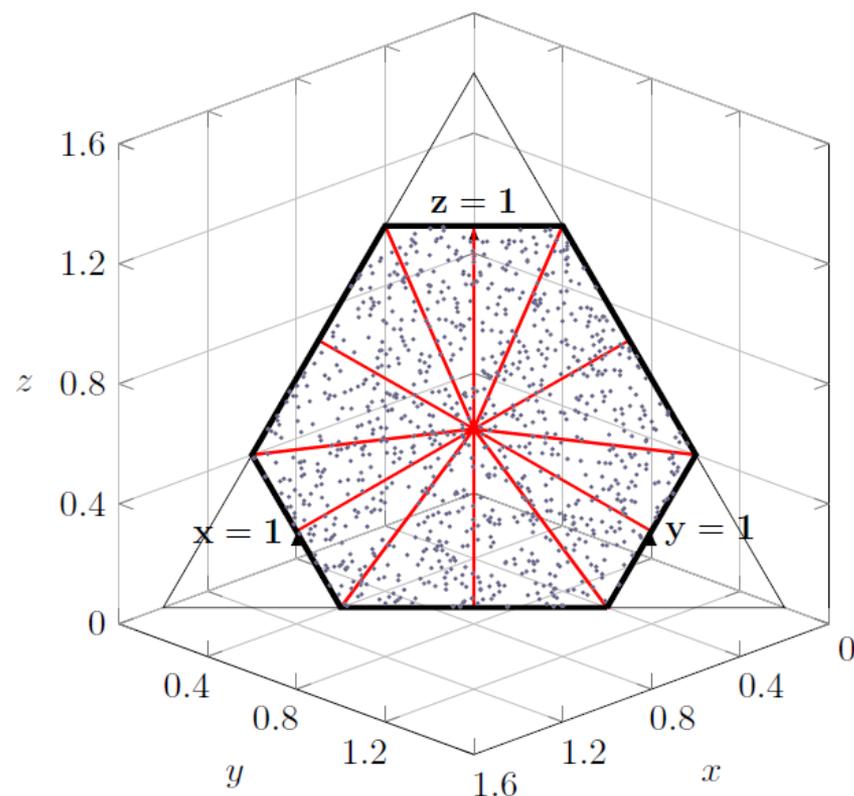
- Davis and Burns (2010) [14]
- Developed for multiprocessor systems, where $U > 1$, but $U_i > 1$ is invalid
- Addresses the problem of maximum (and minimum) constraints
- Very simple (naïve) approach – uses UUnifast then discards any points that do not comply with the constraints
- Suffers from the *curse of dimensionality*: If the constraints on each component halve the volume of the valid region then the proportion of useful points is $1/2^n$ (fine when $n=3$, not so good when $n=50$)



Related work

■ RandFixedSum

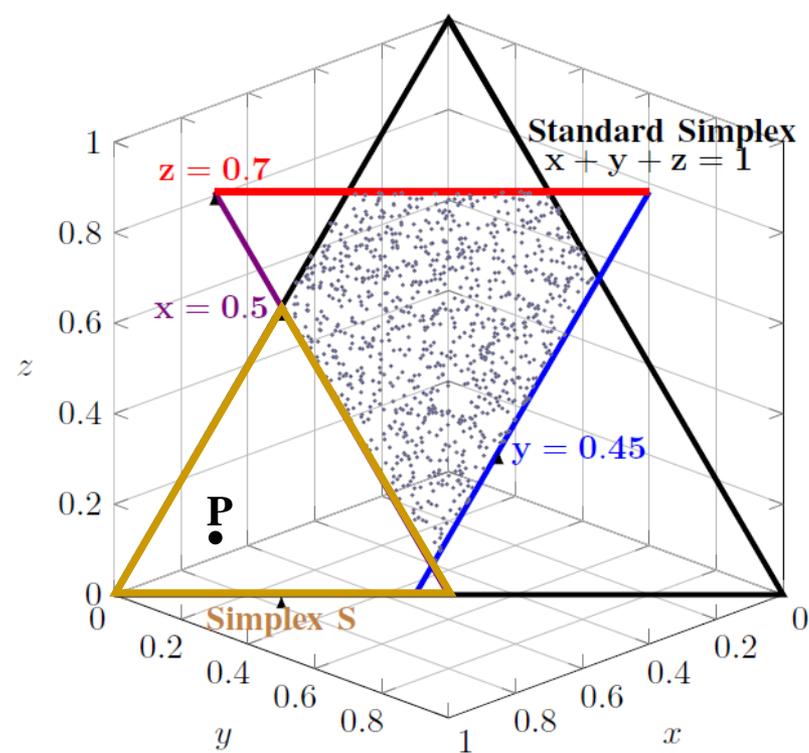
- Invented by Stafford, 2006 [28] and adapted for task set generation by Emberson et al., 2010 [17]
- Efficiently addresses the problem of *symmetric* maximum and minimum constraints (i.e. the same constraints for all tasks)
- De facto standard approach for modelling multiprocessor systems
- Does not cater for *asymmetric* constraints and cannot be adapted to do so because of its reliance on symmetry for its efficiency



Dirichlet-Rescale (DRS) algorithm

■ DRS algorithm

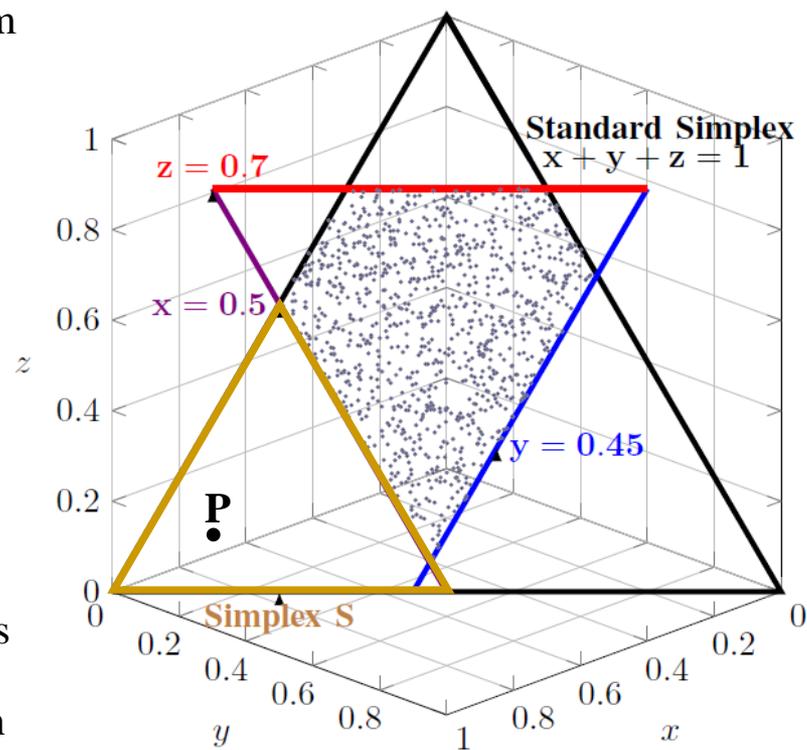
- Addresses the intractability drawbacks: of UUnifast-Discard (discarding points) and of RandFixedSum (would need to generate points in very many different simplices to deal with a valid region that is an irregular shape)
- Basic concept is to generate a point in the standard simplex then if it is not in the valid region, make a series of transformations shifting the coordinates of the point until it is within the valid region
- Crucially these transformations must preserve the uniform distribution of points



How DRS works

■ DRS Algorithm outline operation

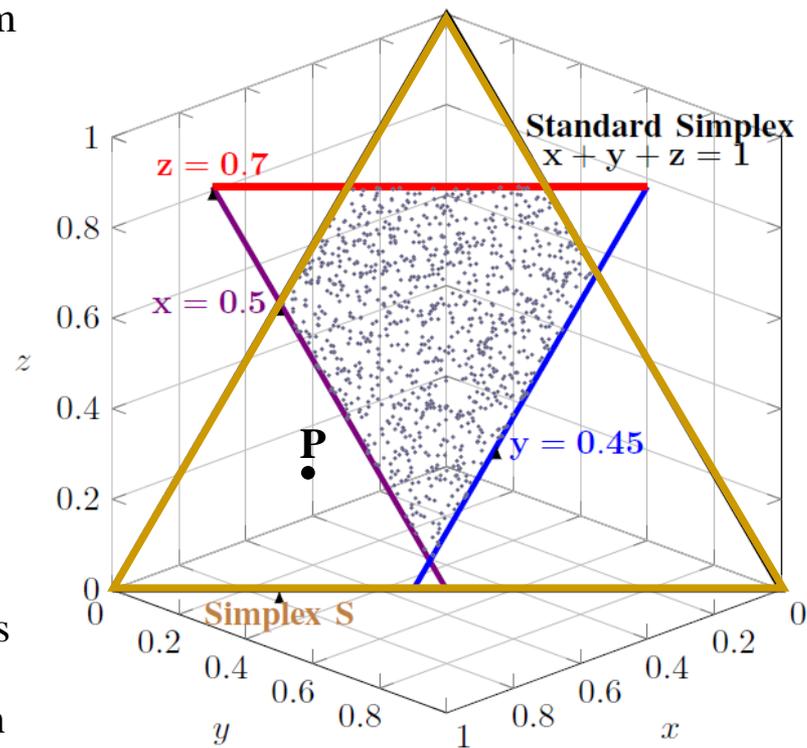
1. Transform the problem into a canonical form by removing minimum constraints
2. Exploits duality to switch the standard and constraints simplices for efficiency
3. Generate a point \mathbf{P} on the standard simplex using the Dirichlet distribution
4. If \mathbf{P} satisfies the constraints then return \mathbf{P} (reversing the initial transformation)
5. Otherwise, defines **Simplex S** based on the broken constraints (**S** contains \mathbf{P})
6. Map **Simplex S** onto the standard simplex via a matrix transformation
7. This scale and translate transformation alters the coordinates of \mathbf{P} making it more likely that the point will now be in the valid region
8. Goto step 4.



How DRS works

■ DRS Algorithm outline operation

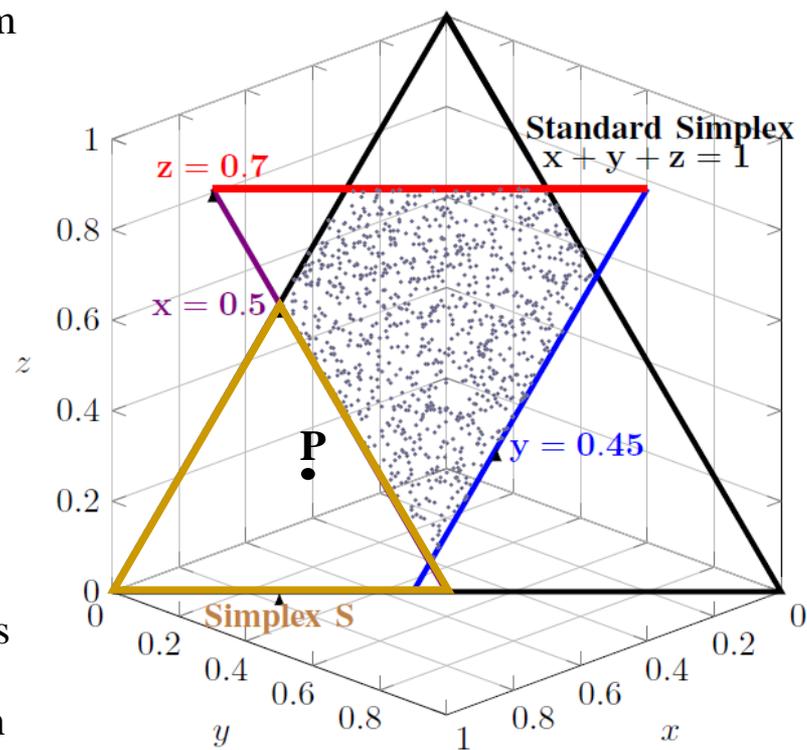
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How DRS works

■ DRS Algorithm outline operation

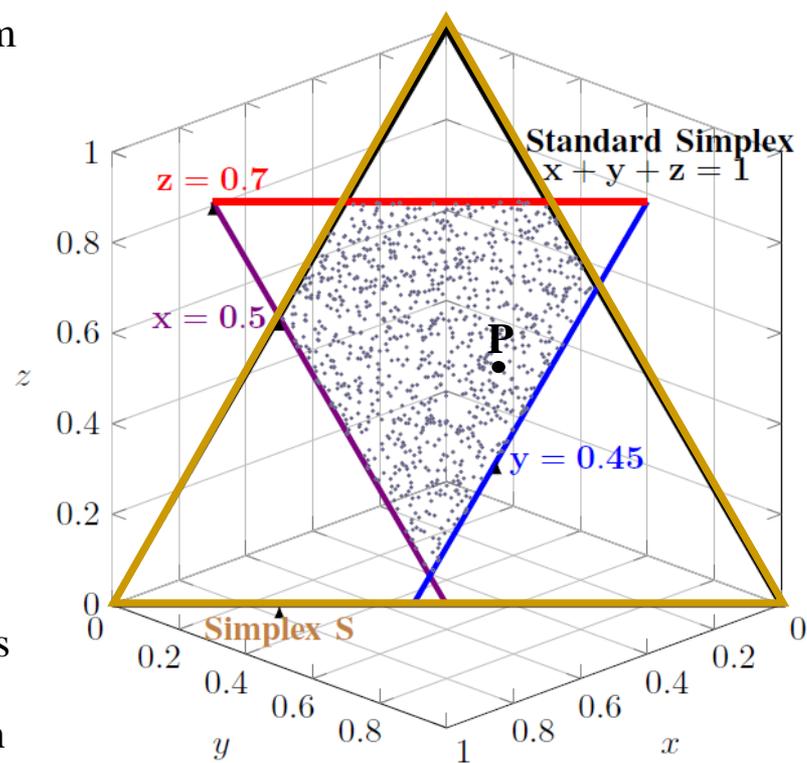
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How DRS works

■ DRS Algorithm outline operation

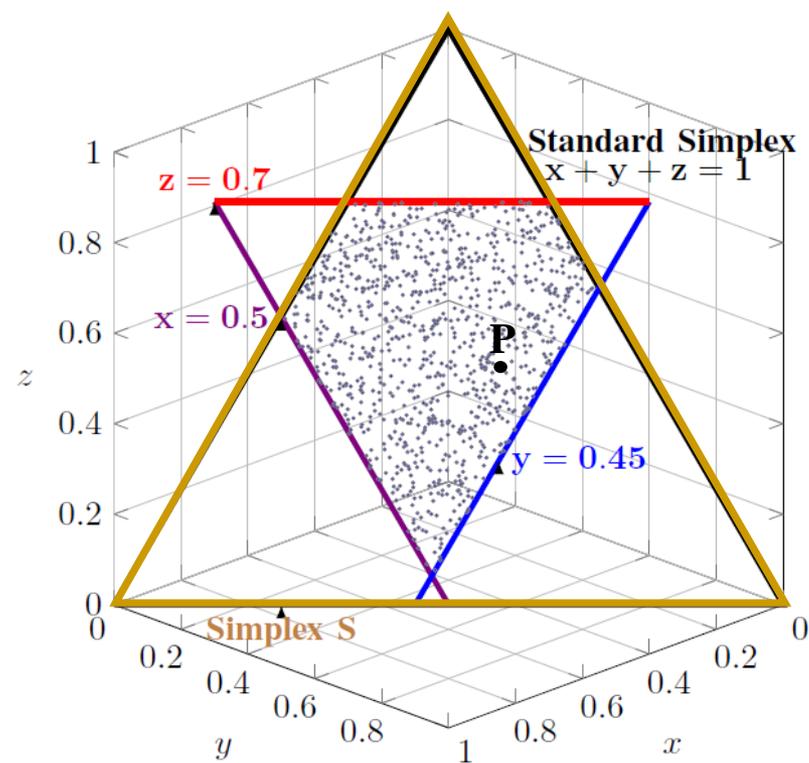
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How DRS works

■ Ensuring Uniformity

- Distribution of initial points generated over the standard simplex is uniform
- Hence the distribution of points is also uniform over **Simplex S**
- The matrix transformation that maps **Simplex S** onto the standard simplex is an *Affine* transformation (i.e. a scale and translate transformation).
- Therefore the points that are uniformly distributed over **Simplex S** become uniformly distributed over the standard simplex and hence uniformly distributed over the valid region



How DRS works (convergence)

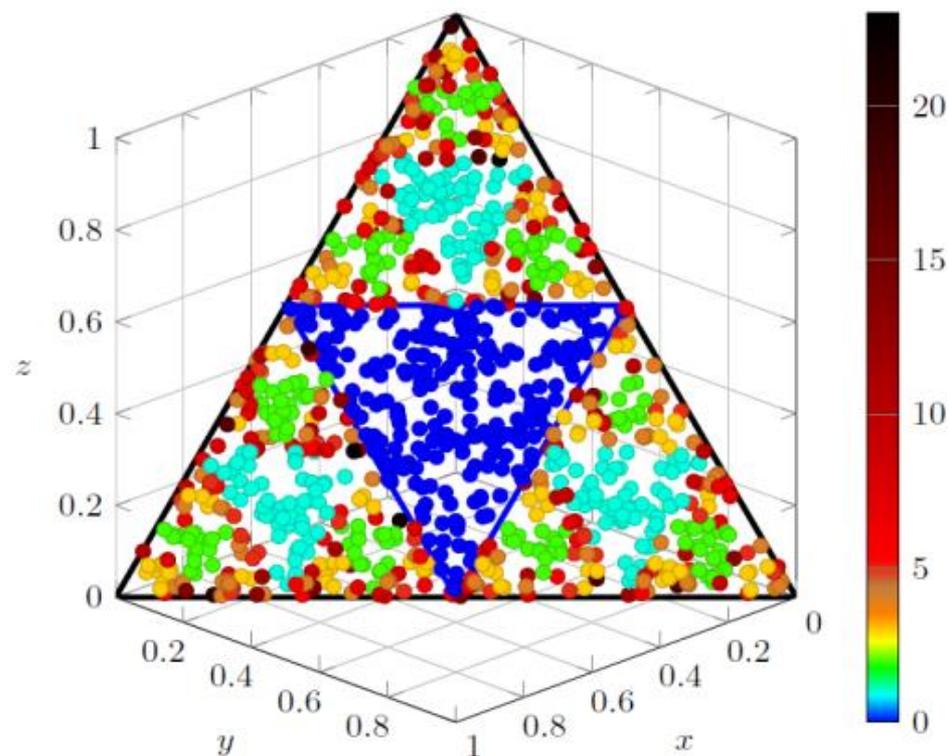
■ Convergence

- Let $p = \frac{\text{volume}(\text{valid region})}{\text{volume}(\text{standard simplex})}$
- After q iterations, the minimum converged volume $c \geq 1 - (1 - p)^q$
- As $q \rightarrow \infty$, $c \rightarrow 1$ and so the algorithm converges

■ Illustration of convergence

- Heat map color codes the number of rescales needed to converge:
- Here $p = 0.25$ and all 1000 initial points converged within 24 rescales

[Note this was done for illustration purposes with the duality optimization disabled, otherwise no rescaling would be necessary since every point generated would be within the smaller constraints simplex (blue)]



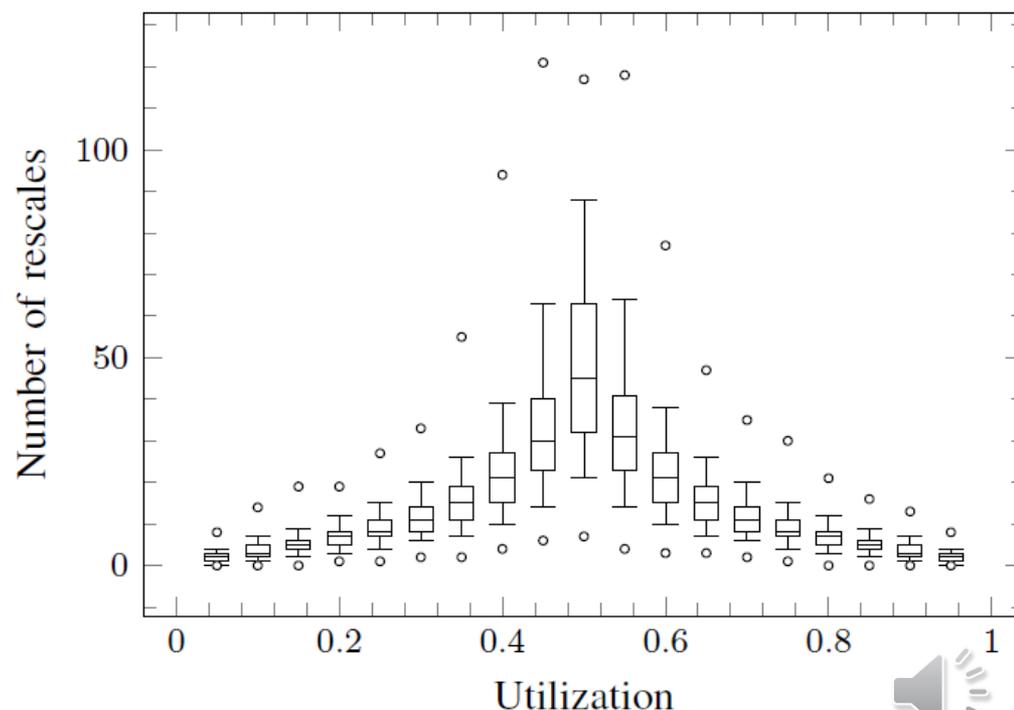
DRS Performance

■ Experiment A

- $n = 50$ and 10,000 runs for each U in $[0.05, 0.95]$ in steps of 0.05
- For each run: $\text{DRS}(n, U, \mathbf{u}^{\max})$ with constraints $\mathbf{u}^{\max} = \text{UUnifast}(n, 1)$

■ Number of Rescales (Box plot)

- Worst-case occurs for $U = 0.5$ when constraints and standard simplex are the same size
- Max rescales < 200 (upper circle)
Min rescales (lower circle)
Mean (middle line of box)
Percentiles (5%, 25%, 75%, 95%)



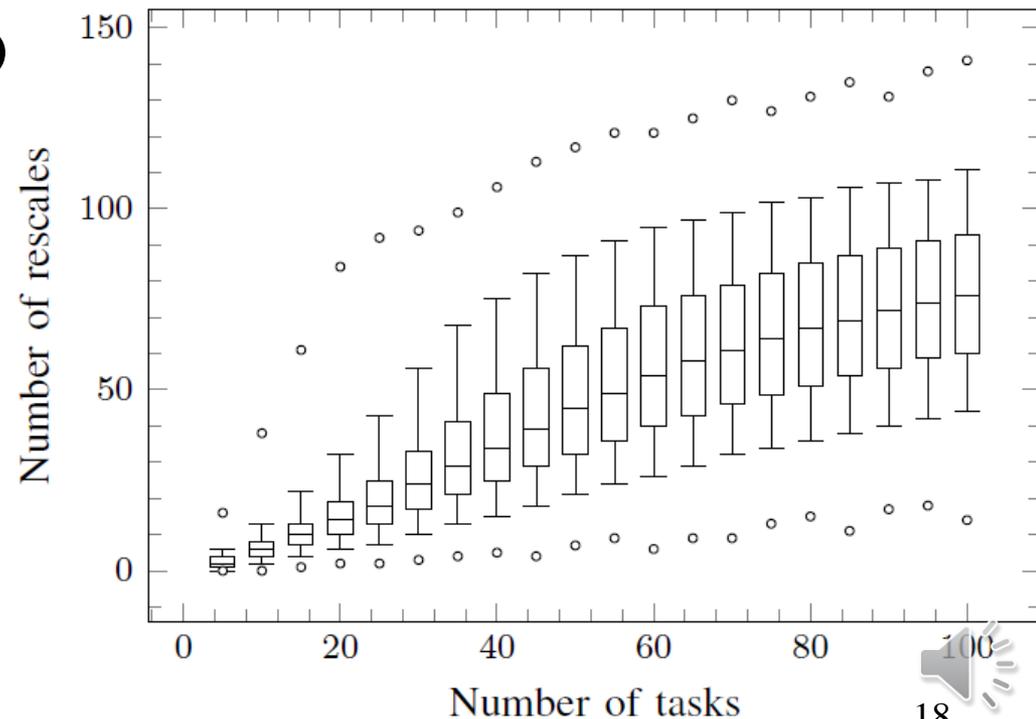
DRS Performance

■ Experiment B

- Similar to Expt. A, but U fixed at 0.5 and n varied from 5 to 100 in steps of 5
- For each run: $\text{DRS}(n, U, \mathbf{u}^{\max})$ with constraints $\mathbf{u}^{\max} = \text{UUnifast}(n, 1)$

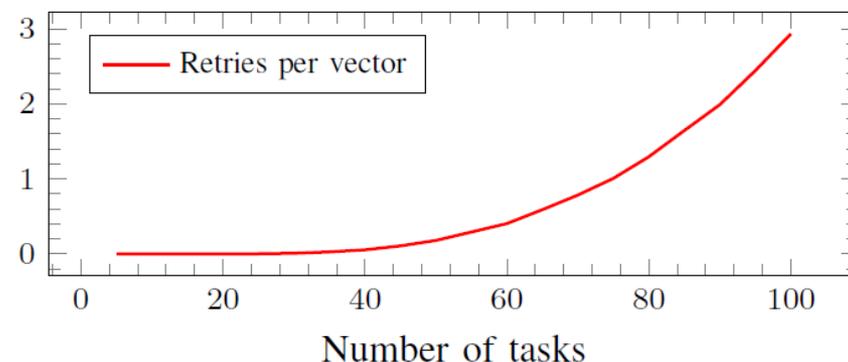
■ Number of Rescales (Box plot)

- Number of rescales gradually increases with increasing size of the vectors (number of tasks)
- Max rescales < 200 (upper circle)
Min rescales (lower circle)
Mean (middle line of box)
Percentiles (5%, 25%, 75%, 95%)



DRS Performance

- **Experiment B (continued)**
- **Number of Retries**
 - Rescale operations can lead to the accumulation of Floating Point error
 - A retry is done by generating another point if the total error (sum of component values minus required utilization) exceeds 0.01%
 - Number of retries increases with increasing size of the vectors, but remains low for $n \leq 100$



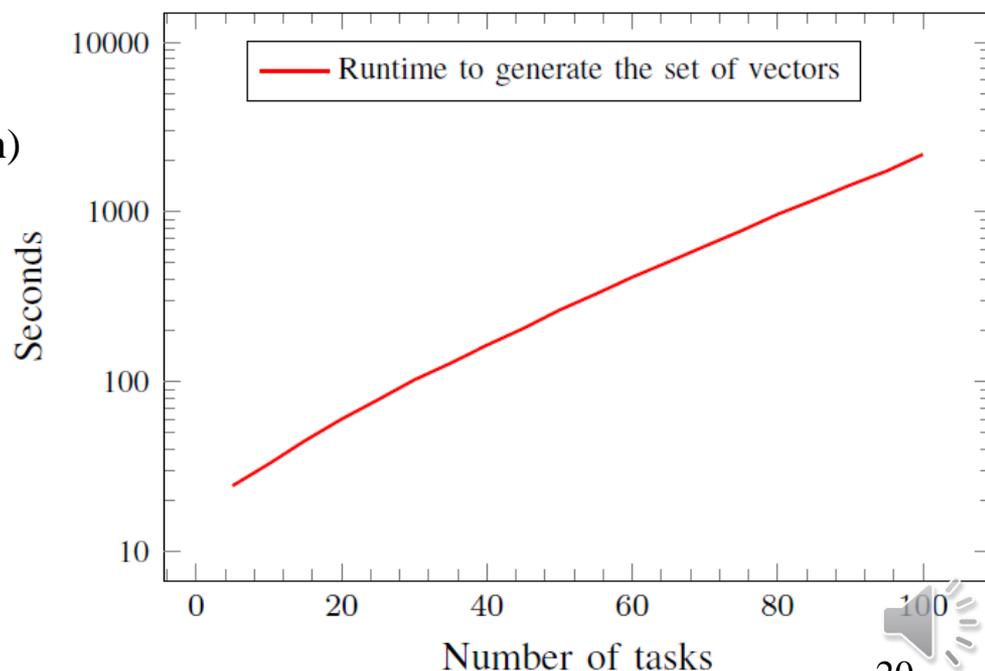
DRS Performance

■ Experiment C

- Runtime to generate all the vectors needed for a standard “benchmark” schedulability analysis experiment (1000 vectors for each of 18 utilization levels from $U = 0.05$ to 0.95 in steps of 0.05, 18,000 vectors in all)

■ Runtimes:

- Used a Pi 4 to obtain reliable timings
- Runtimes well approximated by a polynomial of order 3 (cubic function) $R^2 = 0.999$
- Typical use would be on a laptop or desktop PC (e.g. Dell XPS 13 with Intel™ i7-1065G7 at 3.5GHz
 - ~6 seconds for 10 tasks
 - ~60 seconds for 50 tasks
 - ~6 minutes for 100 tasks(approx. 6 times faster than a Pi 4)



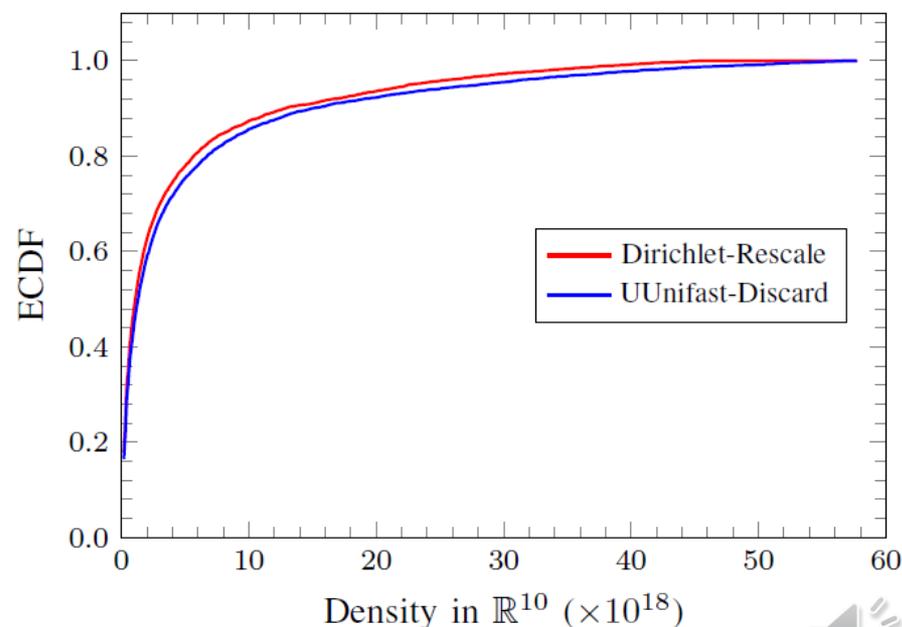
DRS Performance

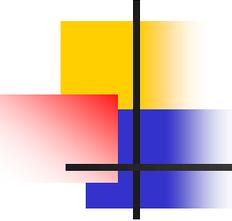
■ Experiment D

- Verified the uniformity of the distribution of vectors produced by DRS via comparison with UUnifast-Discard

■ Statistical test:

- Examined the density of points produced in 1000 small reference simplices (within the valid region) via the DRS algorithm and UUnifast-Discard
- Compared the Empirical Cumulative Distribution Functions (ECDF) using a statistical test: Kolmogorov-Smirnov (KS) test
- KS-statistic = 0.04, p-value = 1.0
- No evidence that the vectors produced come from different distributions
- Cannot reject the null hypothesis that the distributions are the same





Use of the DRS algorithm

- **Main use is in the systematic evaluation of schedulability tests**
 - Used to underpin the generation of synthetic task sets with execution times derived from the utilization values
- **Asymmetric constraints:**
 - Occur when execution times have multiple values or are composed from multiple parts:
 - Mixed Criticality Systems (e.g. C(LO), C(HI))
 - Multi-core systems (e.g. processor demand, bus demand, memory demand, etc.),
 - Typical and worst-case execution times
 - Self-suspensions and resource locking
- **No constraints or symmetric constraints:**
 - DRS can be used to replace UUnifast for single processor systems, and RandFixedSum and UUnifast-Discard for multiprocessor systems

Mixed Criticality Systems Example

■ **Schedulability Analysis Experiment**

- Reproduced from the Adaptive Mixed Criticality (AMC) scheduling paper (Baruah et al., 2011 [4])

■ **Using DRS:**

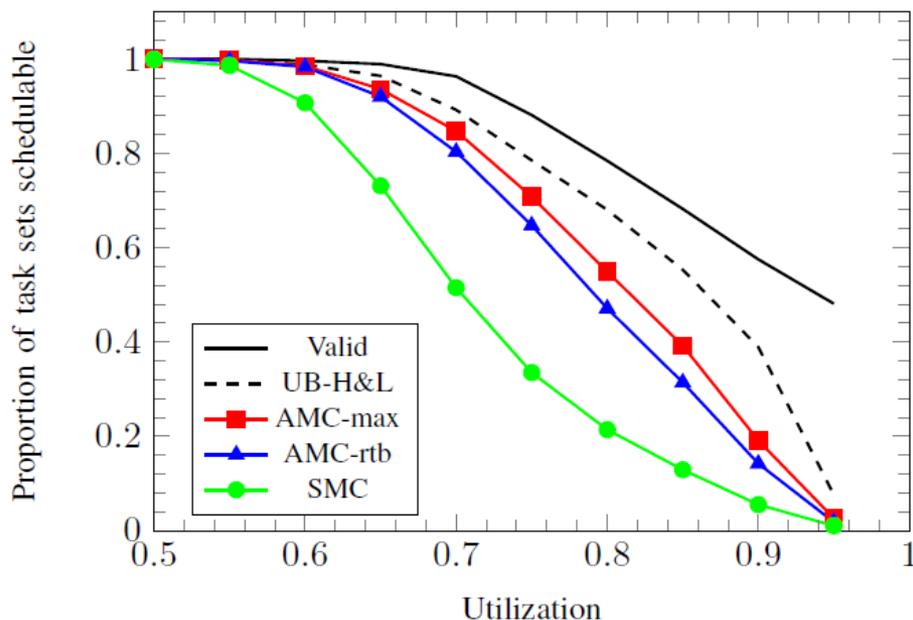
- Independent control of total $U(LO)$ and $U(HI)$
- Independent selection of $U_i(LO) \leq U_i(HI)$ and hence $C_i(LO) \leq C_i(HI)$
- Eliminates generation of invalid (infeasible) task sets
- $U_i(HI)$ generated by calling $\text{DRS}(n^{HI}, U_{HI}^{HI}, \mathbf{u}^1)$
- Maximum constraints set to 1 for LO-criticality tasks and to $U_i(HI)$ for HI-criticality tasks
- $U_i(LO)$ generated by calling $\text{DRS}(n, U^{LO}, \mathbf{u}^{\max})$

Mixed Criticality Systems Example

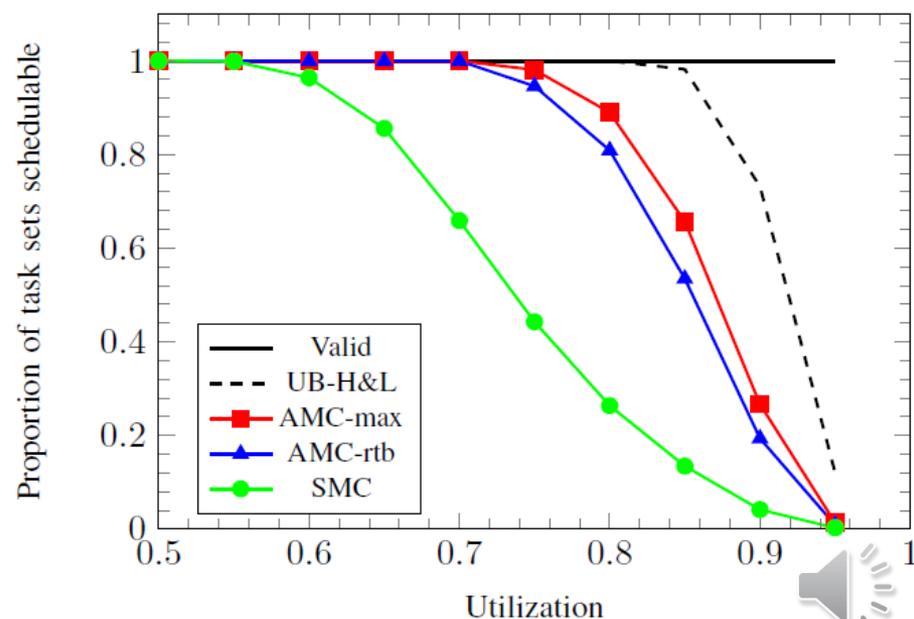
■ Schedulability Analysis Experiment

- Reproduced from AMC paper [4]
- DRS highlights sharper transition of AMC and larger improvement over SMC
- More nuanced and realistic results – could affect decisions on which methods to use

Baruah et al.



DRS used in task set generation



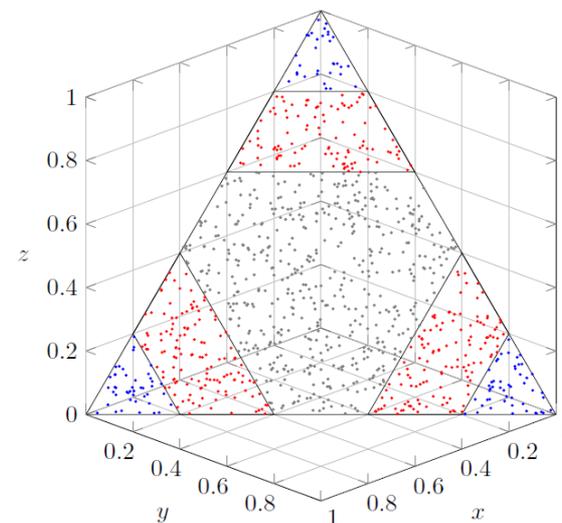
Why use an unbiased distribution of utilization vectors?

■ What is meant by an unbiased?

- Vectors generated are uniformly distributed across the valid region
- Does not mean the component values themselves are uniformly distributed (common misconception)

■ Why use a unbiased distribution?

- For generic schedulability analysis experiments, using a uniform distribution of utilization vectors means that each possible vector that complies with the constraints has the same chance of being selected
- The distribution is thus unbiased, provides full and fair coverage of all valid possibilities, and is therefore arguably the appropriate one to use
- Not using a uniform distribution of vectors risks biasing the results of schedulability analysis experiments



Easy ways of introducing bias...

1. Confound variables (n and U)

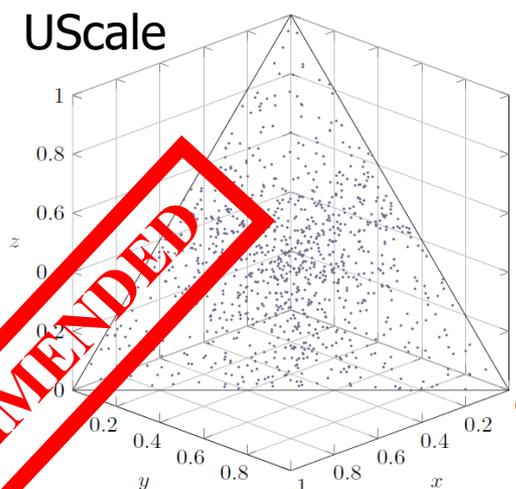
- Select U_i from a uniform distribution $[0,1]$ and keep adding tasks until the required total utilization is reached
- Confounds n and U , so we cannot distinguish the effects of higher task set cardinality from those of higher task set utilization

2. Simple scaling (UScale)

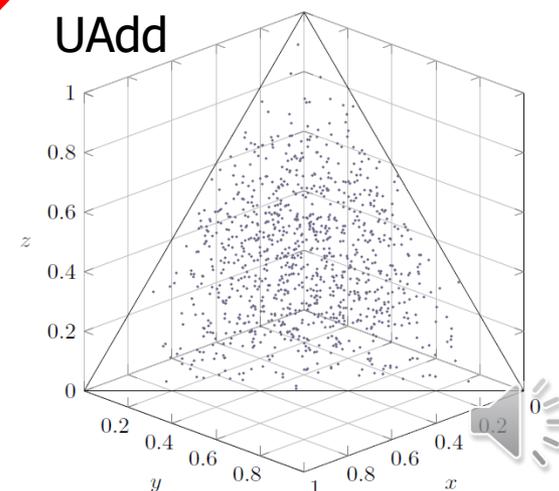
- Select n values for U_i from a uniform distribution $[0,1]$ and then scale them to achieve the required total utilization U

3. Addition of components (UAdd)

- Use UUnifast for each of multiple parts of U_i and then add these values together



NOT RECOMMENDED



Conclusion:

Why use the DRS algorithm?



■ Flexible - general purpose algorithm

- Supports asymmetric constraints on maximum and minimum utilization for each task
 - Used to obtain unbiased distributions when execution times have multiple values or are composed from multiple parts
 - Useful for tailoring task sets to specific problem requirements, limitations, or domain specific constraints
- Can also be used to replace UUnifast, UUnifast-Discard, and RandFixedSum

■ High performance

- Supports efficient generation of task sets with cardinality up to $n = 100$ with individual constraints
- Additional experiments show that DRS supports generation of task sets with cardinality up to $n = 200$ with a commensurate slowdown in performance

■ Python source code is publicly available

- Permanently archived at <https://doi.org/10.5281/zenodo.4118059>
- Can be installed via: `pip install drs` (<https://pypi.org/project/drs/>)
- Also provide a C library enabling the DRS algorithm in Python to be called directly from C/C++ code



Questions?

